

# Side-Channel and Fault Attacks in Modern Lattice-Based Cryptography

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Julius Hermelink

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It has been more than 10 years since I started at Infineon 😊

- 2014-2020: Studying mathematics.
- 2014: Started at Infineon at DES.
- 2018: Switched to CCS/Thomas Pöppelmann.
- 2020: Begin of PhD in cooperation with UniBW M.
- 2023: Started at MPI-SP in Bochum
- 2024: Finished PhD



# The Quantum Threat

Quantum computers threaten currently used asymmetric cryptography.



We have to assume that:

- Large-scale quantum computer break commonly used asymmetric schemes.
- Adversaries: harvest now, decrypt later.

Therefore, we need:

- Post-quantum asymmetric cryptography.
- Most pressingly key exchanges.

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# The NIST Standardization Process

NIST is in the process of standardizing post-quantum cryptography.

<sup>1</sup> [FIPS 203 \(Draft\)](#)

<sup>2</sup> Federal Information Processing Standards Publication  
<sup>3</sup>

<sup>4</sup> **Module-Lattice-based  
Key-Encapsulation  
Mechanism Standard**

<sup>7</sup> Category: Computer Security      Subcategory: Cryptography

<sup>8</sup> Information Technology Laboratory  
<sup>9</sup> National Institute of Standards and Technology  
<sup>10</sup> Gaithersburg, MD 20899-8900

<sup>11</sup> This publication is available free of charge from:  
<sup>12</sup> <https://doi.org/10.6023/NIST.FIPS.203.ipd>

<sup>13</sup> Published August 24, 2023



<sup>14</sup>

<sup>15</sup> U.S. Department of Commerce  
<sup>16</sup> Gina M. Raimondo, Secretary

<sup>17</sup> National Institute of Standards and Technology

<sup>18</sup> Laurie E. Locascio, NIST Director and Under Secretary of Commerce for Standards and Technology

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- Four candidates already selected.
- Three are lattice-based.
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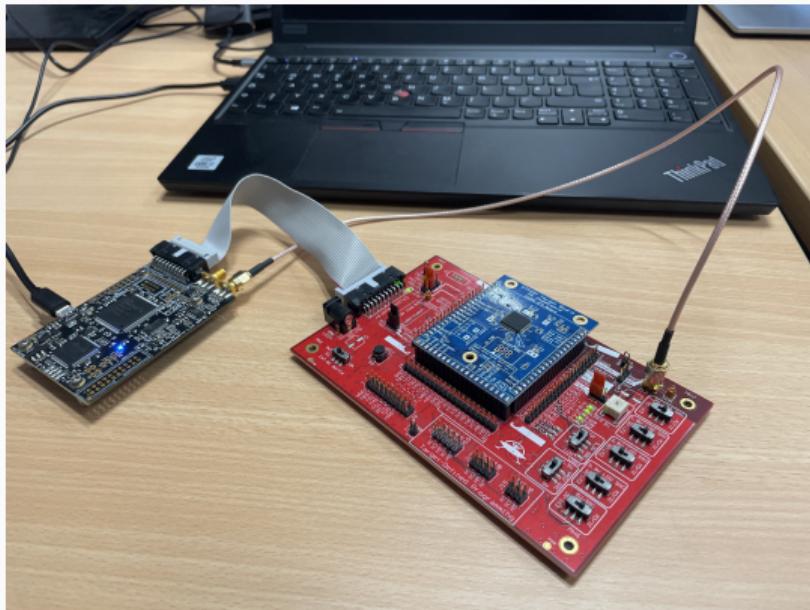
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ML-KEM used in Signal, Chrome, iMessage, ...



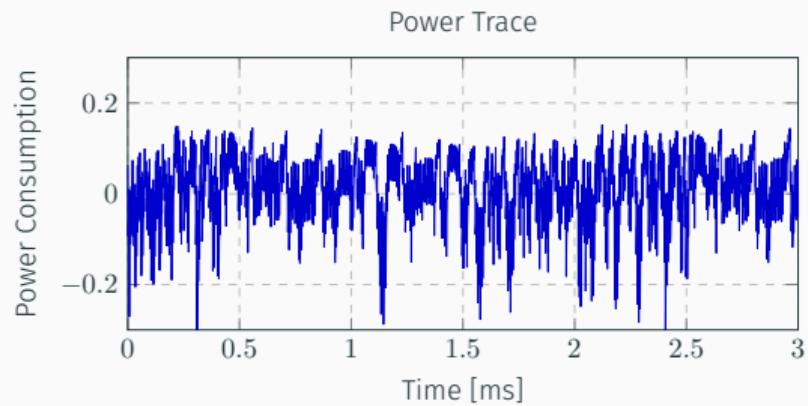
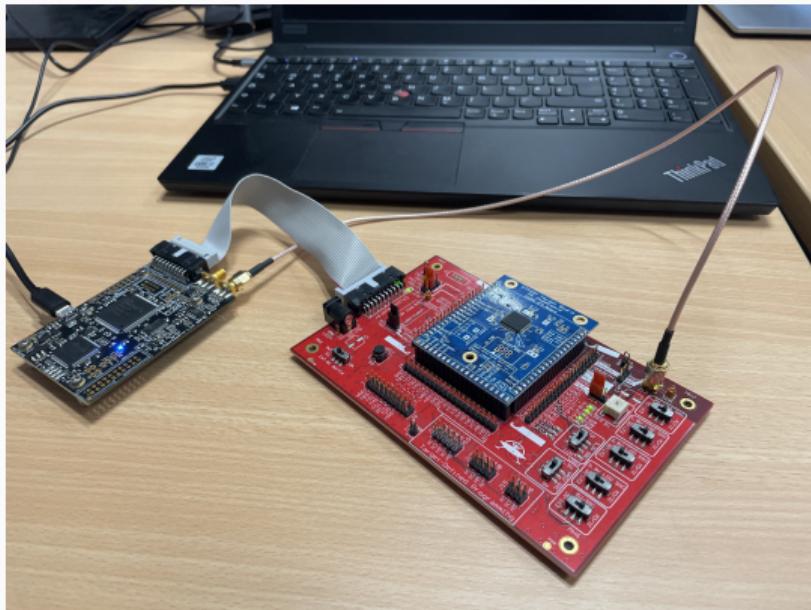
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- Different underlying hard problems.
- Different multiplications (e.g., using number theoretic transforms).
- Error correction to recover message from noisy coefficients.
- Construction from PKE using FO-transforms to achieve IND-CCA security.
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The number theoretic transform (NTT):

- Enables fast multiplication in several lattice-based schemes.
- Used at multiple points in all routines of ML-KEM.
- Inverse NTT processes data depending on the secret key during decryption.
- Previous work established (inverse) NTT as target for side-channel attacks.
- However, required noise levels limit attacks when targeting secret key.

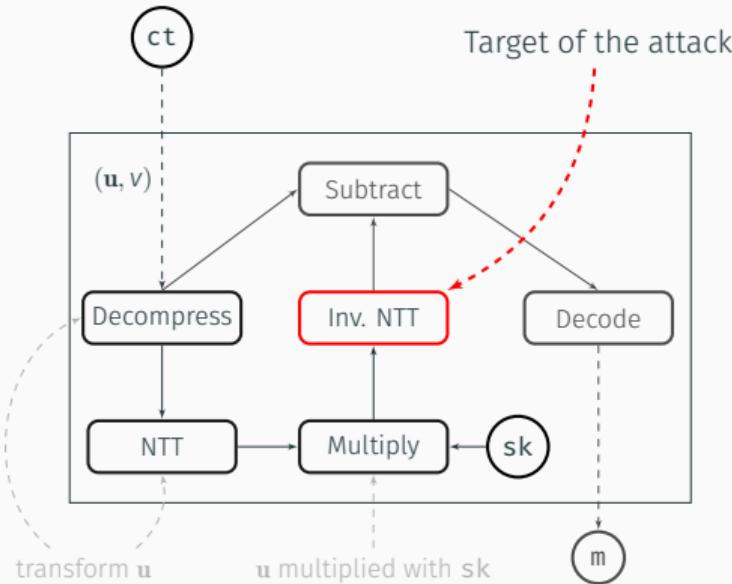
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# Chosen-Ciphertext k-Trace Attacks – Idea



Previous work [PPM17, PP19]:

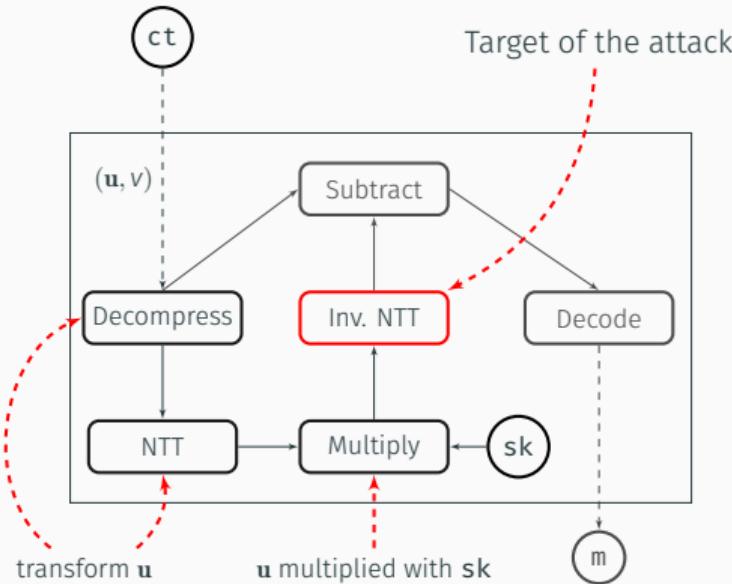
- Template attack on inv. NTT; then belief propagation.
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Decryption as shown on the left:

- Ciphertext components are decompressed.
- Component is multiplied with secret.
- Results fed into the inv. number theoretic transform.

Attack strategy: Reduce entropy using compressible NTT-sparse chosen ciphertext.

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# Chosen-Ciphertext k-Trace Attacks

Our attack strategy for increased noise tolerance:

For targeted subkeys:

- Formulate as lattice problem.
- Run lattice reduction.
- Obtain compressible NTT-sparse ct.

For each ct:

- Record trace for ct.
- Obtain distributions for intermediates.
- Run belief propagation; obtain subkeys.

Using the subkeys:

- Formulate key recovery using subkeys as lattice problem.
- Run lattice reduction.
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**Lattice reduction is computationally expensive and slow but done offline.**

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# Adapting Belief Propagation to Counter Shuffling of NTTs

Real-world attacks have to take countermeasures into account.

Important classes of countermeasures:

- Masking circumvented.
- Hiding prevents these attacks.

Ravi et al. [RPBC20] (ascending security):

- Fine shuffling
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- Fine shuffling: Shuffle node adapts factor depending on processed information.
- Coarse shuffling: Extended attacker model and matching algorithm.

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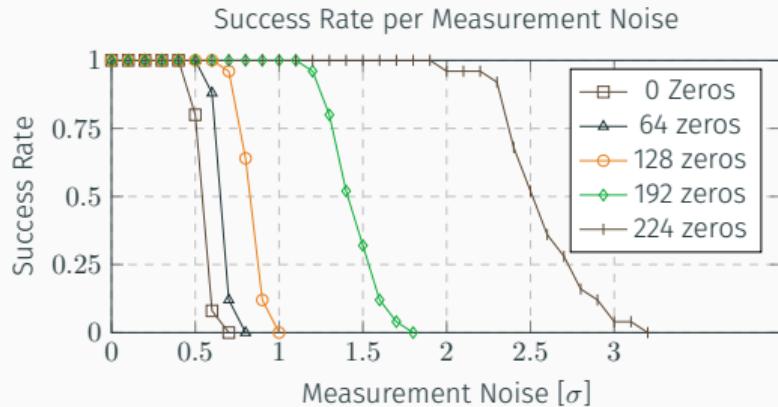
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Evaluation in the leakage models provided by previous work [PPM17, PP19].



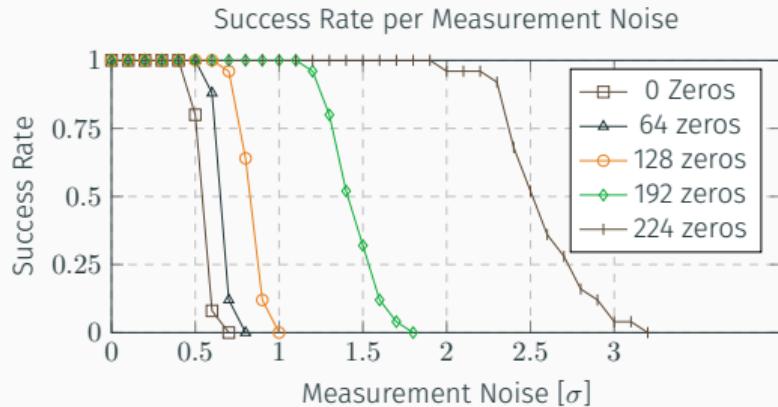
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This work	Yes	$\sigma \leq 1.7$ (3.1)	1-8	Yes

Noise tolerance increased from  $\sigma \leq 0.6$  to  $\sigma \leq 1.7$  ( $\sigma \leq 3.1$ ).

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# Targeting the FO-Transform

Learning with errors schemes recover the message from noisy coefficients:

- Observation whether introduced error causes failure: leaks information.
- Fujisaki-Okamoto (FO) transform achieves IND-CCA2 security.

Previous attacks:

- CCA to potentially cause failure; observe using SCA on comparison [GJN20, BDH<sup>+</sup>21].
- Or use fault against decoder to potentially cause failure and observe outcome [PP21].
- Require insufficiently protected comparison/decoder; reliable fault.

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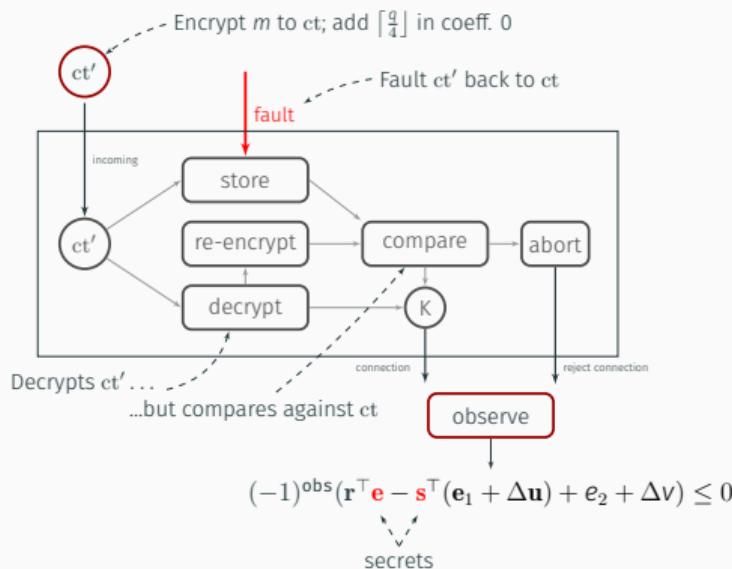
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# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

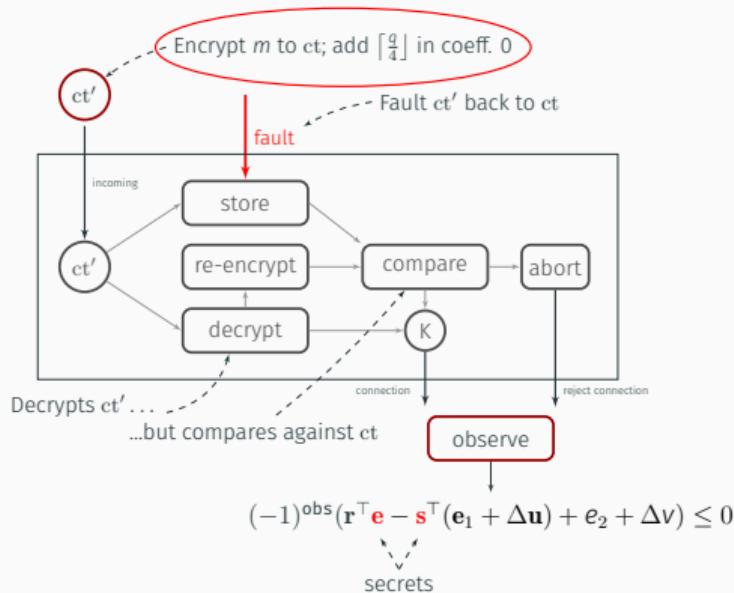


Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.

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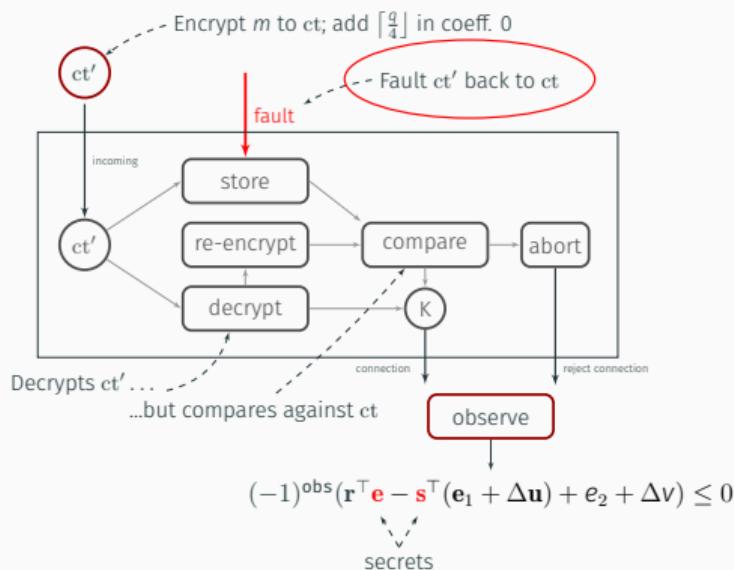


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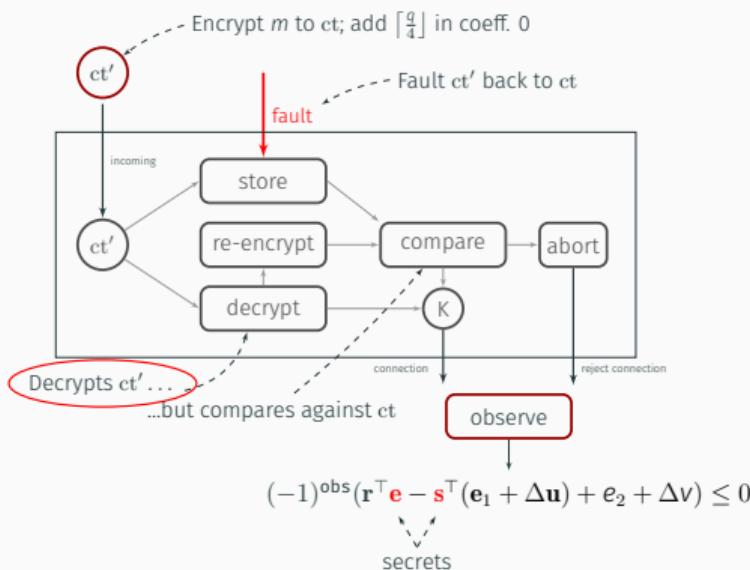


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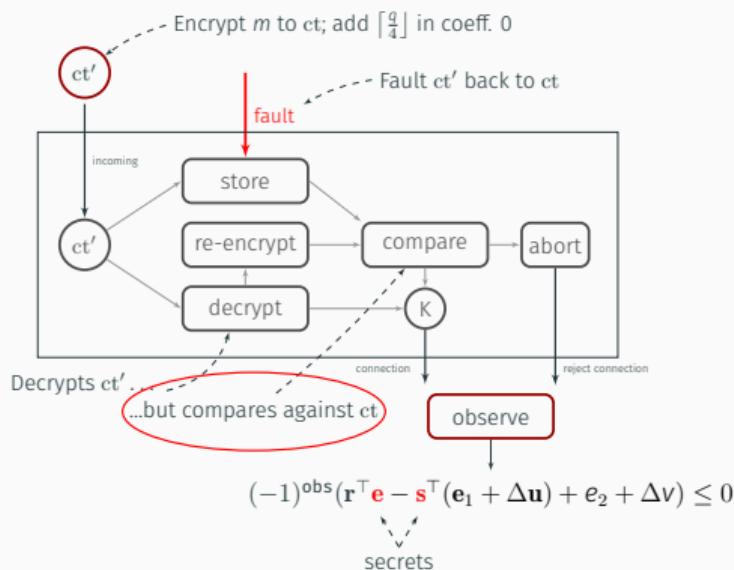


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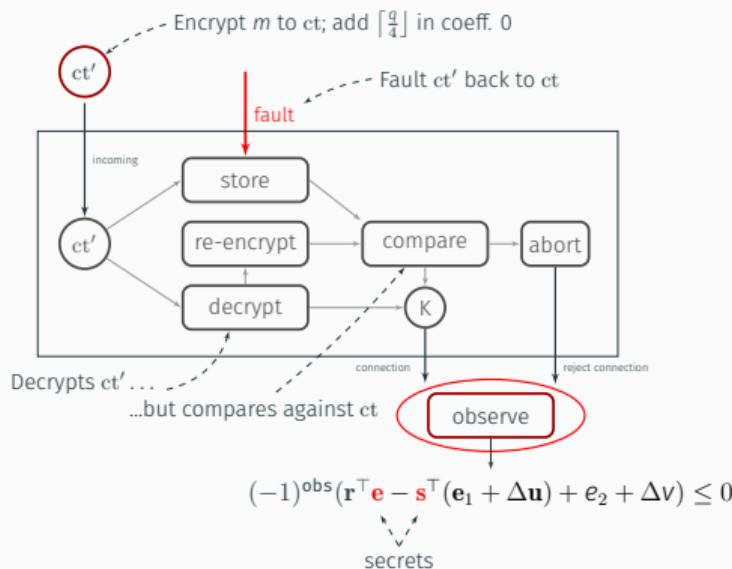


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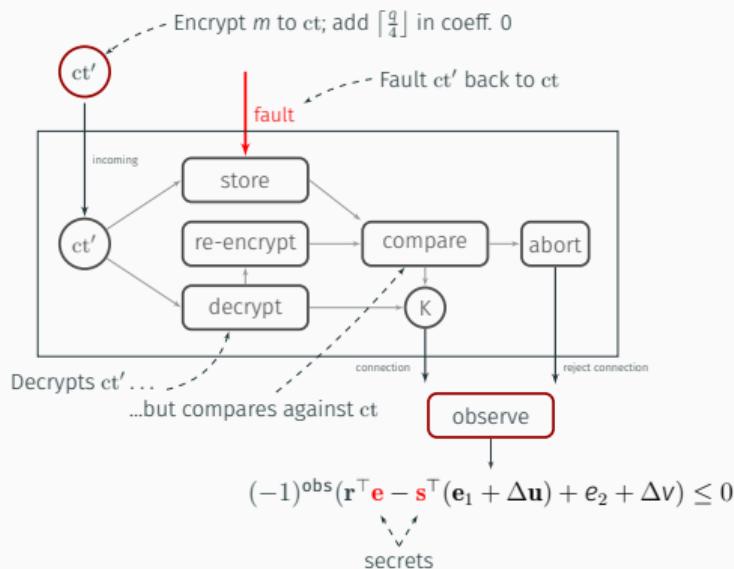


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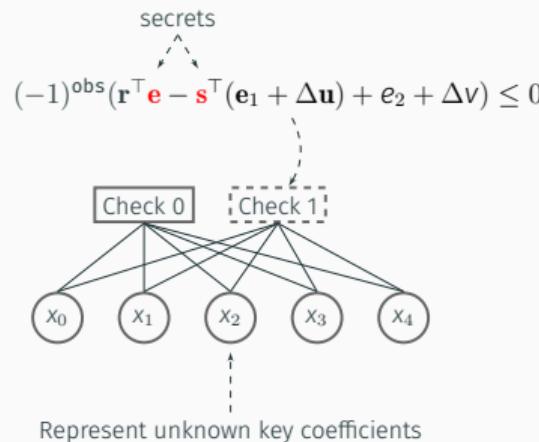


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# Fault-Enabled Chosen-Ciphertext Attacks – BP

We propose solving decryption failure inequalities using belief propagation.



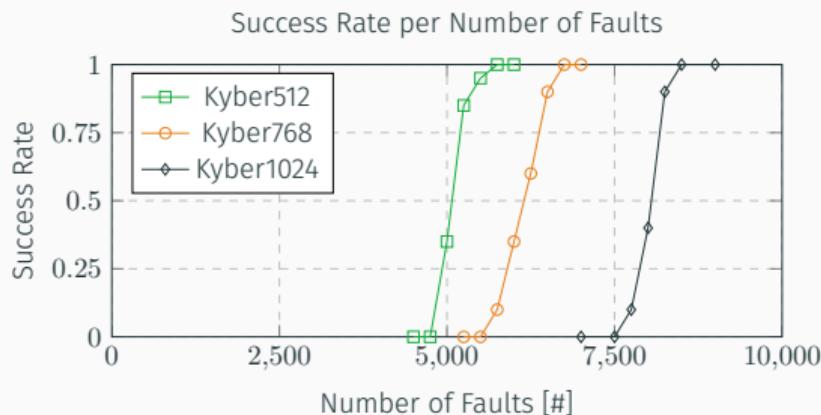
Belief propagation inspired by [PP21]:

- Check nodes represent inequalities.
- Variable nodes represent unknown coefficients.
- Priors are binomial distributions of secrets.

Requires fewer inequalities while being more computationally efficient.

# Fault-Enabled Chosen-Ciphertext Attacks – Results

Gives a more general class of attacks resulting from our method.

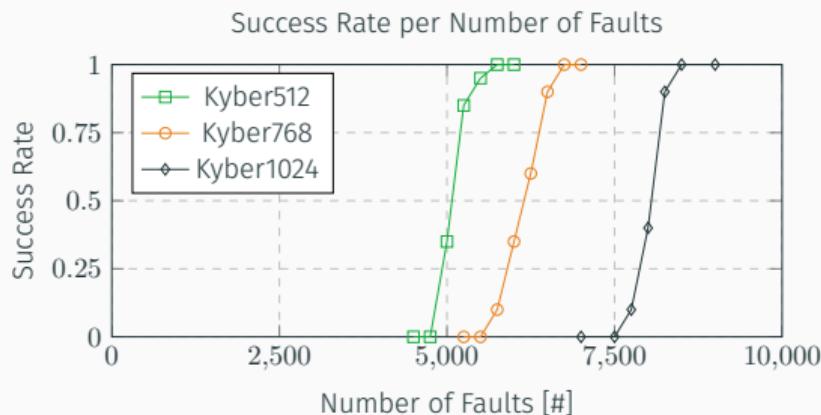


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Decryption failures leak information in form of inequalities:

- Exploited in wide variety of attacks including ours.
- Particularly hard to mitigate.
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Attacks require recovery method to solve for secret key:

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- Error resistance increases #inequalities.
- General problem: combine belief propagation and algebraic methods?

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# Security Estimates for Error-Tolerant Key Recovery

Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

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How can we combine the advantages of previous methods?

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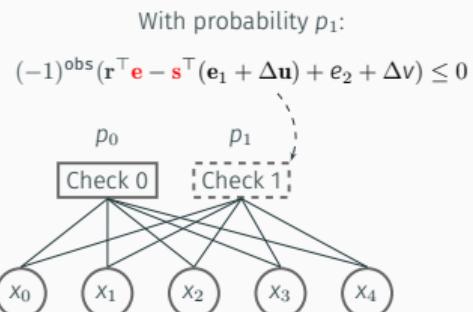
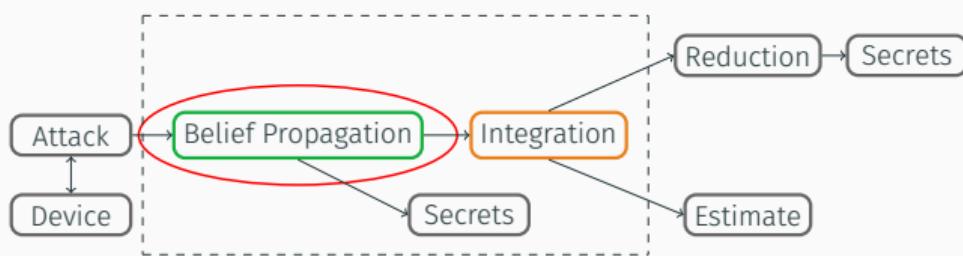
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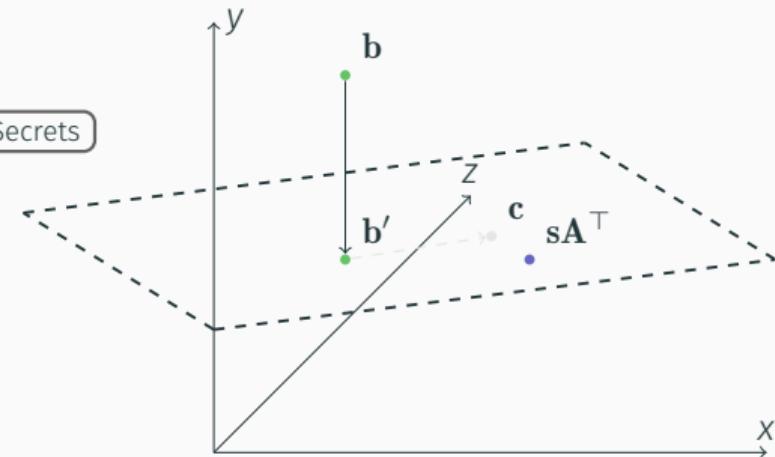
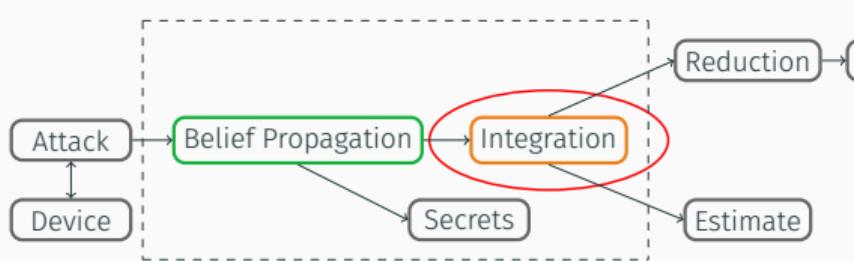
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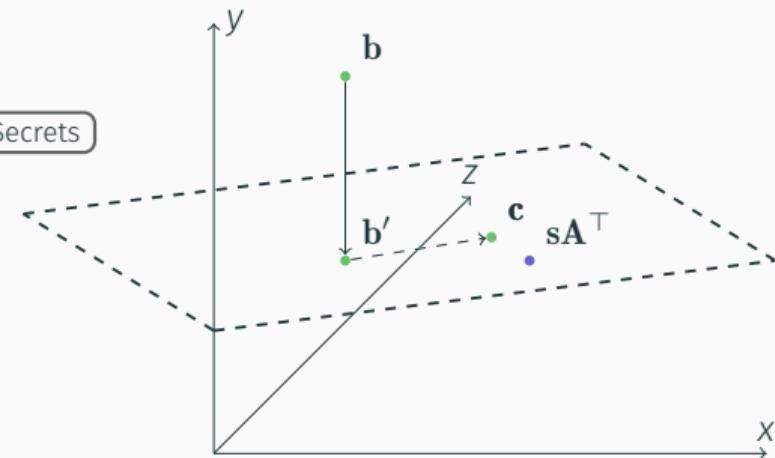
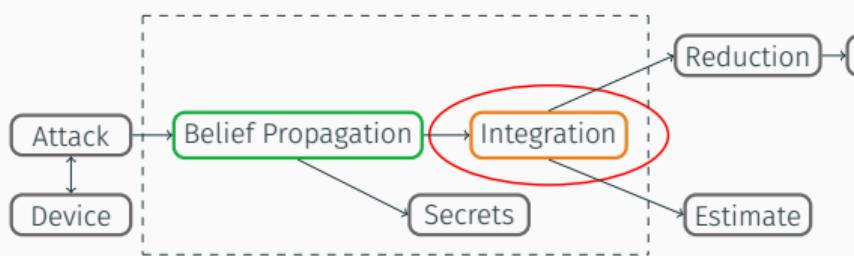
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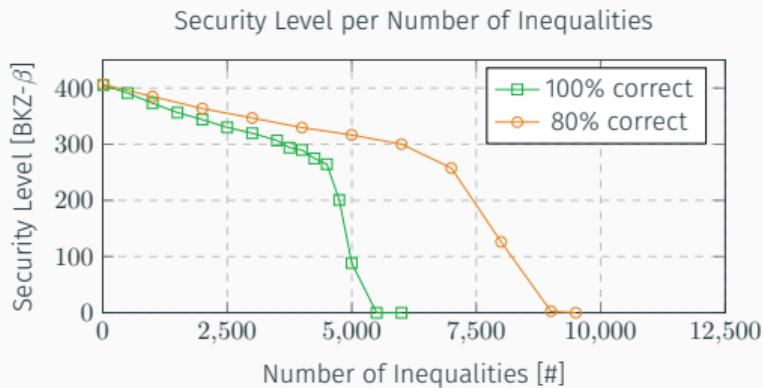


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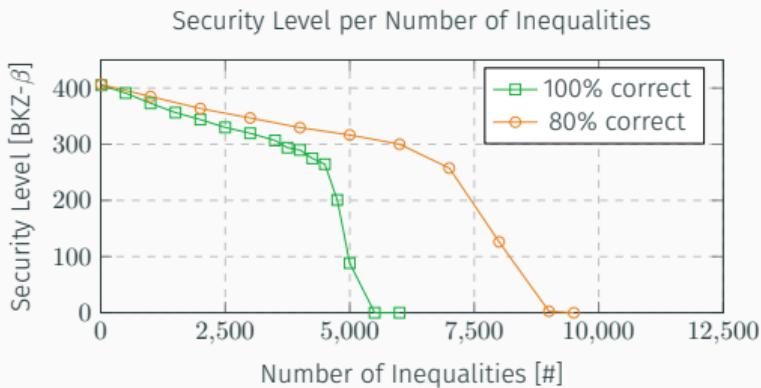
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Hybrid approach: fewer inequalities, error-tolerant, and provides security estimates.

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[Del22]	9000	Yes	No
[DDGR20, DGHK22]	$\geq 10000$	No	Yes
This work [HMS+23b]	5500	Yes	Yes

Hybrid approach: fewer inequalities, error-tolerant, and provides security estimates.

**NIST has standardized ML-KEM – a lattice-based scheme.**

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

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# Publications

[HHP+21] Mike Hamburg, **Julius Hermelink**, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. "Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber". In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.4 (2021), pp. 88–113. URL: <https://doi.org/10.46586/tches.v2021.i4.88-113>

[HSST23] **Julius Hermelink**, Silvan Streit, Emanuele Strieder, and Katharina Thieme. "Adapting Belief Propagation to Counter Shuffling of NTTs". In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.1 (2023), pp. 60–88. URL: <https://doi.org/10.46586/tches.v2023.i1.60-88>

[HMS+23b] **Julius Hermelink**, Erik Mårtensson, Simona Samardjiska, Peter Pessl, and Gabi Dreßler. "Belief Propagation Meets Lattice Reduction: Security Estimates for Error-Tolerant Key Recovery from Decryption Errors". In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.4 (2023), pp. 287–317. URL: <https://doi.org/10.46586/tches.v2023.i4.287-317>

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[HPP21] **Julius Hermelink**, Peter Pessl, and Thomas Pöppelmann. "Fault-Enabled Chosen-Ciphertext Attacks on Kyber". In: *Progress in Cryptology - INDOCRYPT 2021 - 22nd International Conference on Cryptology in India, Jaipur, India, December 12-15, 2021, Proceedings*. Ed. by Avishek Adhikari, Ralf Küsters, and Bart Preneel. Vol. 13143. Lecture Notes in Computer Science. Springer, 2021, pp. 311–334. URL: [https://doi.org/10.1007/978-3-030-92518-5\\_15](https://doi.org/10.1007/978-3-030-92518-5_15)

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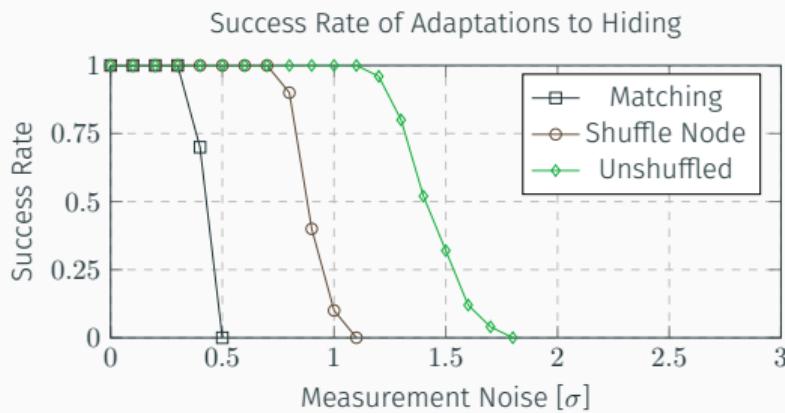
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# Adapting Belief Propagation – Results

Real-world attacks have to take countermeasures into account.



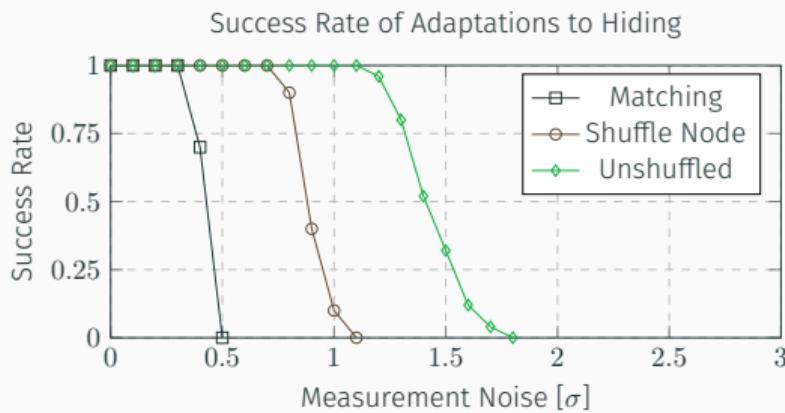
Our results show that

- Attacks not fully prevented by hiding countermeasures.
- However, noise tolerance reduced.
- Strongest form of shuffling protects requires vast computational resources.
- Large-scale adversaries might even circumvent coarse full shuffling.

Several hiding countermeasures can be circumvented; protected NTT still vulnerable.

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