

# Side-Channel and Fault Attacks in Modern Lattice-Based Cryptography

---

Julius Hermelink

Max Planck Institute for Security and Privacy

# Years at Infineon

It has been more than 10 years since I started at Infineon 😊

- 2014-2020: Studying mathematics.
- 2014: Started at Infineon at DES.
- 2018: Switched to CCS/Thomas Pöppelmann.
- 2020: Begin of PhD in cooperation with UniBW M.
- 2023: Started at MPI-SP in Bochum
- 2024: Finished PhD



# The Quantum Threat

Quantum computers threaten currently used asymmetric cryptography.



We have to assume that:

- Large-scale quantum computer break commonly used asymmetric schemes.
- Adversaries: harvest now, decrypt later.

Therefore, we need:

- Post-quantum asymmetric cryptography.
- Most pressing key exchanges.

# The Quantum Threat

Quantum computers threaten currently used asymmetric cryptography.



We have to assume that:

- Large-scale quantum computer break commonly used asymmetric schemes.
- Adversaries: harvest now, decrypt later.

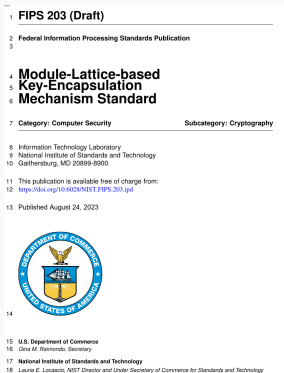
Therefore, we need:

- Post-quantum asymmetric cryptography.
- Most pressing key exchanges.



# The NIST Standardization Process

NIST is in the process of standardizing post-quantum cryptography.

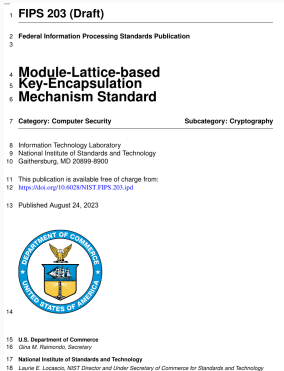


NIST started a standardization process in 2016.

- Fourth round ongoing.
- Four candidates already selected.
- Three are lattice-based.
- Kyber selected as KEM (Kyber  $\mapsto$  ML-KEM).

# The NIST Standardization Process

NIST is in the process of standardizing post-quantum cryptography.



NIST started a standardization process in 2016.

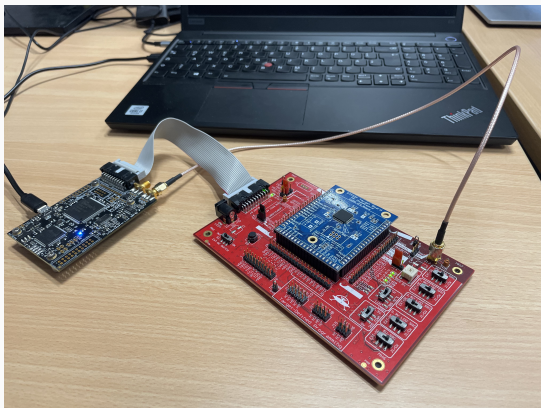
- Fourth round ongoing.
- Four candidates already selected.
- Three are lattice-based.
- Kyber selected as KEM (Kyber  $\mapsto$  ML-KEM).

ML-KEM used in Signal, Chrome, iMessage, ...



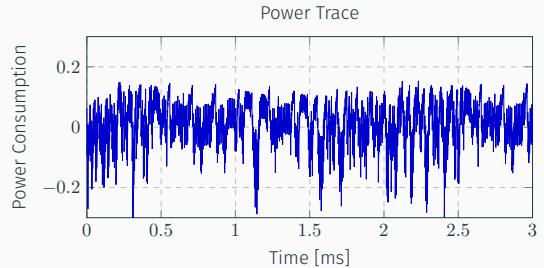
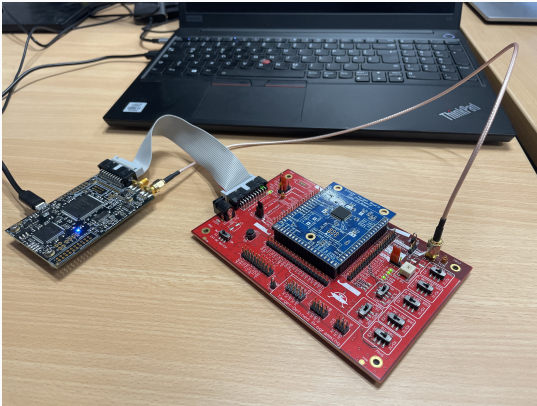
# Attacks on Embedded Devices

Embedded devices may be vulnerable to side-channel and fault attacks.



# Attacks on Embedded Devices

Embedded devices may be vulnerable to side-channel and fault attacks.



Lattice-Based Cryptography uses different building blocks.

- Different underlying hard problems.
- Different multiplications (e.g., using number theoretic transforms).
- Error correction to recover message from noisy coefficients.
- Construction from PKE using FO-transforms to achieve IND-CCA security.
- ...

Which new vulnerabilities in regard to side-channel and fault attacks does this open up?

Lattice-Based Cryptography uses different building blocks.

- Different underlying hard problems.
- Different multiplications (e.g., using number theoretic transforms).
- Error correction to recover message from noisy coefficients.
- Construction from PKE using FO-transforms to achieve IND-CCA security.
- ...

Which new vulnerabilities in regard to side-channel and fault attacks does this open up?

The number theoretic transform (NTT):

- Enables fast multiplication in several lattice-based schemes.
- Used at multiple points in all routines of ML-KEM.
- Inverse NTT processes data depending on the secret key during decryption.
- Previous work established (inverse) NTT as target for side-channel attacks.
- However, required noise levels limit attacks when targeting secret key.

To what extent is the number theoretic transform vulnerable to side-channel analysis?

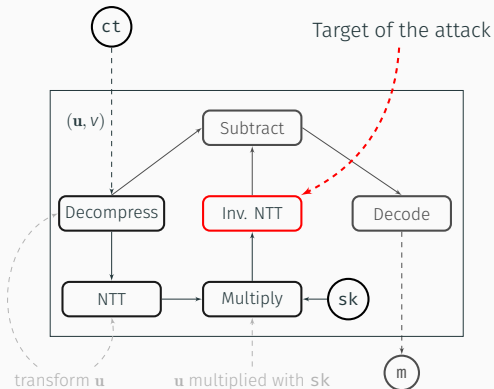
The number theoretic transform (NTT):

- Enables fast multiplication in several lattice-based schemes.
- Used at multiple points in all routines of ML-KEM.
- Inverse NTT processes data depending on the secret key during decryption.
- Previous work established (inverse) NTT as target for side-channel attacks.
- However, required noise levels limit attacks when targeting secret key.

**To what extent is the number theoretic transform vulnerable to side-channel analysis?**



# Chosen-Ciphertext k-Trace Attacks – Idea



Previous work [PPM17, PP19]:

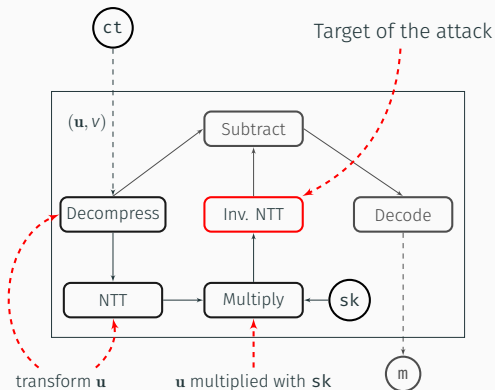
- Template attack on inv. NTT; then belief propagation.
- However, cannot target secret key with high noise tolerance.

Decryption as shown on the left:

- Ciphertext components are decompressed.
- Component is multiplied with secret.
- Results fed into the inv. number theoretic transform.

Attack strategy: Reduce entropy using compressible NTT-sparse chosen ciphertext.

# Chosen-Ciphertext k-Trace Attacks – Idea



Previous work [PPM17, PP19]:

- Template attack on inv. NTT; then belief propagation.
- However, cannot target secret key with high noise tolerance.

Decryption as shown on the left:

- Ciphertext components are decompressed.
- Component is multiplied with secret.
- Results fed into the inv. number theoretic transform.

Attack strategy: Reduce entropy using compressible NTT-sparse chosen ciphertext.

# Chosen-Ciphertext k-Trace Attacks

Our attack strategy for increased noise tolerance:

For targeted subkeys:

- Formulate as lattice problem.
- Run lattice reduction.
- Obtain compressible NTT-sparse  $\mathbf{ct}$ .

For each  $\mathbf{ct}$ :

- Record trace for  $\mathbf{ct}$ .
- Obtain distributions for intermediates.
- Run belief propagation; obtain subkeys.

Using the subkeys:

- Formulate key recovery using subkeys as lattice problem.
- Run lattice reduction.
- Obtain full key.

**Lattice reduction is computationally expensive and slow but done offline.**

# Chosen-Ciphertext k-Trace Attacks

Our attack strategy for increased noise tolerance:

For targeted subkeys:

- Formulate as lattice problem.
- Run lattice reduction.
- Obtain compressible NTT-sparse **ct**.

For each **ct**:

- Record trace for **ct**.
- Obtain distributions for intermediates.
- Run belief propagation; obtain subkeys.

Using the subkeys:

- Formulate key recovery using subkeys as lattice problem.
- Run lattice reduction.
- Obtain full key.

**Lattice reduction is computationally expensive and slow but done offline.**

# Chosen-Ciphertext k-Trace Attacks

Our attack strategy for increased noise tolerance:

For targeted subkeys:

- Formulate as lattice problem.
- Run lattice reduction.
- Obtain compressible NTT-sparse  $\mathbf{ct}$ .

For each  $\mathbf{ct}$ :

- Record trace for  $\mathbf{ct}$ .
- Obtain distributions for intermediates.
- Run belief propagation; obtain subkeys.

Using the subkeys:

- Formulate key recovery using subkeys as lattice problem.
- Run lattice reduction.
- Obtain full key.

**Lattice reduction is computationally expensive and slow but done offline.**

# Adapting Belief Propagation to Counter Shuffling of NTTs

Real-world attacks have to take countermeasures into account.

Important classes of countermeasures:

- Masking circumvented.
- Hiding prevents these attacks.

Ravi et al. [RPBC20] (ascending security):

- Fine shuffling
- Coarse block shuffling
- Coarse full shuffling

We propose two techniques against hiding:

- Fine shuffling: Shuffle node adapts factor depending on processed information.
- Coarse shuffling: Extended attacker model and matching algorithm.

Adaptation to hiding countermeasures for belief-propagation-based attacks.

# Adapting Belief Propagation to Counter Shuffling of NTTs

Real-world attacks have to take countermeasures into account.

Important classes of countermeasures:

- Masking circumvented.
- Hiding prevents these attacks.

Ravi et al. [RPBC20] (ascending security):

- Fine shuffling
- Coarse block shuffling
- Coarse full shuffling

We propose two techniques against hiding:

- Fine shuffling: Shuffle node adapts factor depending on processed information.
- Coarse shuffling: Extended attacker model and matching algorithm.

Adaptation to hiding countermeasures for belief-propagation-based attacks.

# Adapting Belief Propagation to Counter Shuffling of NTTs

Real-world attacks have to take countermeasures into account.

Important classes of countermeasures:

- Masking circumvented.
- Hiding prevents these attacks.

Ravi et al. [RPBC20] (ascending security):

- Fine shuffling
- Coarse block shuffling
- Coarse full shuffling

We propose two techniques against hiding:

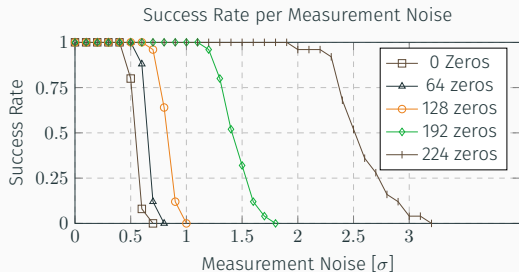
- Fine shuffling: Shuffle node adapts factor depending on processed information.
- Coarse shuffling: Extended attacker model and matching algorithm.

**Adaptation to hiding countermeasures for belief-propagation-based attacks.**



# Chosen-Ciphertext k-Trace Attacks – Results

Evaluation in the leakage models provided by previous work [PPM17, PP19].



Attack	Secret Key	Noise Tolerance	Traces	Hiding considered
[PPM17]	Yes	$\sigma \leq 0.6$	1	No
[PP19]	No	$\sigma \leq 2.0$	1	No
This work	Yes	$\sigma \leq 1.7$ (3.1)	1-8	Yes

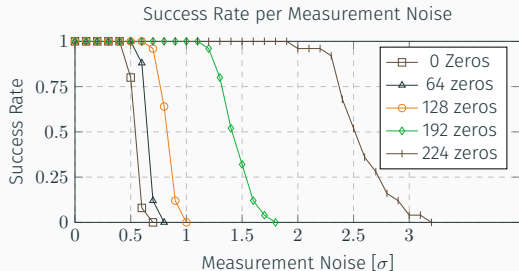
Noise tolerance increased from  $\sigma \leq 0.6$  to  $\sigma \leq 1.7$  ( $\sigma \leq 3.1$ ).

Mike Hamburg, **Julius Hermelink**, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. "Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber". In: *IACR Transactions on Cryptographic Hardware and Embedded Systems 2021.4* (2021), pp. 88–113

**Julius Hermelink**, Silvan Streit, Emanuele Strieder, and Katharina Thieme. "Adapting Belief Propagation to Counter Shuffling of NTTs". In: *IACR Transactions on Cryptographic Hardware and Embedded Systems 2023.1* (2023), pp. 60–88

# Chosen-Ciphertext k-Trace Attacks – Results

Evaluation in the leakage models provided by previous work [PPM17, PP19].



Attack	Secret Key	Noise Tolerance	Traces	Hiding considered
[PPM17]	Yes	$\sigma \leq 0.6$	1	No
[PP19]	No	$\sigma \leq 2.0$	1	No
This work	Yes	$\sigma \leq 1.7$ (3.1)	1-8	Yes

Noise tolerance increased from  $\sigma \leq 0.6$  to  $\sigma \leq 1.7$  ( $\sigma \leq 3.1$ ).

Mike Hamburg, **Julius Hermelink**, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. "Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber". In: *IACR Transactions on Cryptographic Hardware and Embedded Systems 2021.4* (2021), pp. 88–113

**Julius Hermelink**, Silvan Streit, Emanuele Strieder, and Katharina Thieme. "Adapting Belief Propagation to Counter Shuffling of NTTs". In: *IACR Transactions on Cryptographic Hardware and Embedded Systems 2023.1* (2023), pp. 60–88

# Targeting the FO-Transform

Learning with errors schemes recover the message from noisy coefficients:

- Observation whether introduced error causes failure: leaks information.
- Fujisaki-Okamoto (FO) transform achieves IND-CCA2 security.

Previous attacks:

- CCA to potentially cause failure; observe using SCA on comparison [GJN20, BDH<sup>+</sup>21].
- Or use fault against decoder to potentially cause failure and observe outcome [PP21].
- Require insufficiently protected comparison/decoder; reliable fault.

# Targeting the FO-Transform

Learning with errors schemes recover the message from noisy coefficients:

- Observation whether introduced error causes failure: leaks information.
- Fujisaki-Okamoto (FO) transform achieves IND-CCA2 security.

Previous attacks:

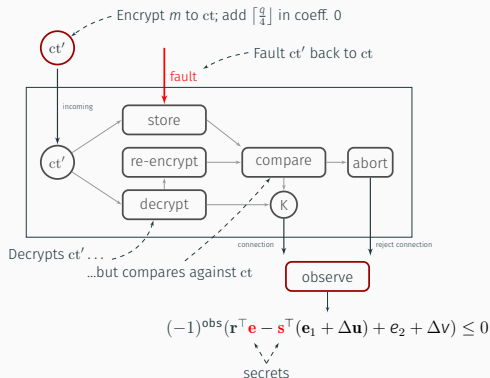
- CCA to potentially cause failure; observe using SCA on comparison [GJN20, BDH<sup>+</sup>21].
- Or use fault against decoder to potentially cause failure and observe outcome [PP21].
- Require insufficiently protected comparison/decoder; reliable fault.

# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

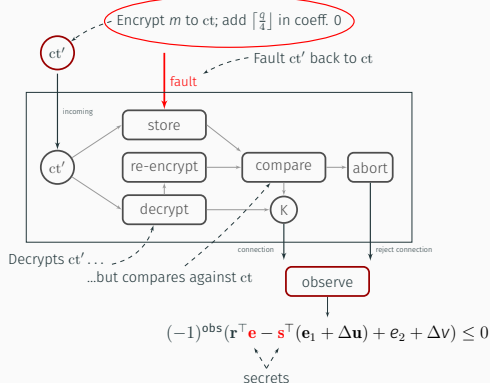
Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.



# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.



Our attack strategy:

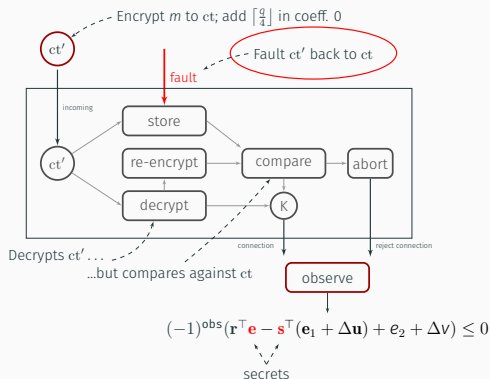
- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.

# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

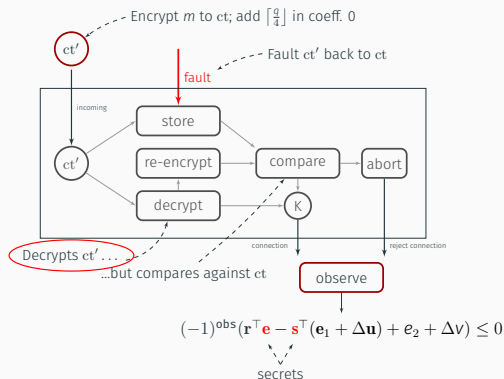
Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.



# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.



Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.

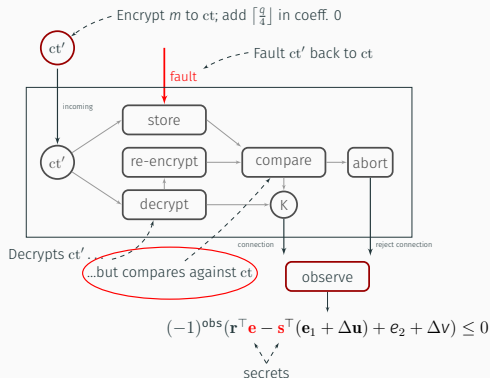


# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.

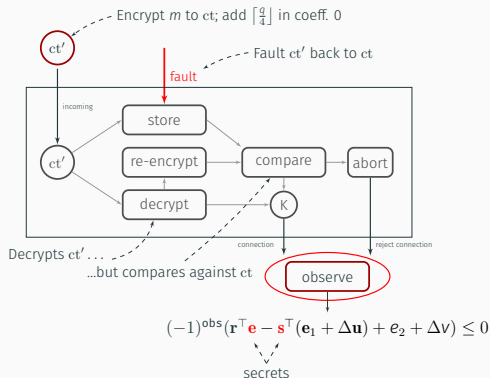


# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.

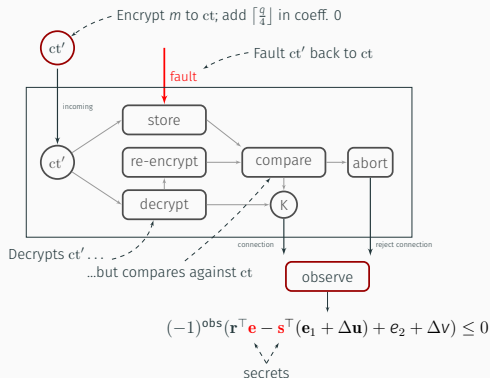


# Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

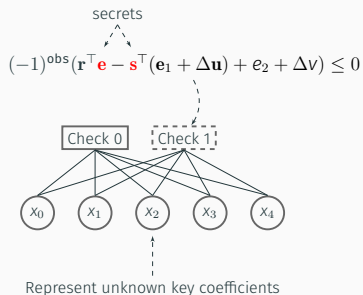
Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts  $ct'$ , but compares to  $ct$ .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.



# Fault-Enabled Chosen-Ciphertext Attacks – BP

We propose solving decryption failure inequalities using belief propagation.



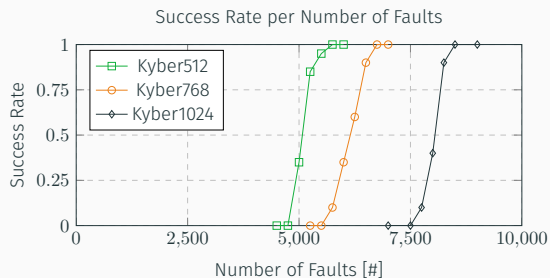
Belief propagation inspired by [PP21]:

- Check nodes represent inequalities.
- Variable nodes represent unknown coefficients.
- Priors are binomial distributions of secrets.

Requires fewer inequalities while being more computationally efficient.

# Fault-Enabled Chosen-Ciphertext Attacks – Results

Gives a more general class of attacks resulting from our method.

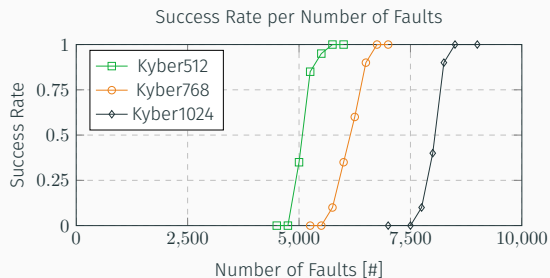


Attack	Type	Point of Attack	Requirement/Robustness
[GJN20]	Timing	Comparison	Non-constant time
[BDH+21]	SCA	Comparison	Leaking comparison
[PP21]	Fault	Decoding	Unprotected decoding
[DHP+22]	SCA	Comparison	Max. first order protection
[De122]	Fault	Multiple	Unreliable/Imprecise fault
[Wei22]	SCA/ML	Multiple	Defeats 1th order masking
[FKK+22]	Rowhammer	Key Generation	KeyGen Failure Boosting
This work [HPP21]	Fault	Multiple	Unreliable fault

Allows for very unreliable fault, enlarges attack surface, only targets public data.

# Fault-Enabled Chosen-Ciphertext Attacks – Results

Gives a more general class of attacks resulting from our method.



Attack	Type	Point of Attack	Requirement/Robustness
[GJN20]	Timing	Comparison	Non-constant time
[BDH+21]	SCA	Comparison	Leaking comparison
[PP21]	Fault	Decoding	Unprotected decoding
[DHP+22]	SCA	Comparison	Max. first order protection
[De22]	Fault	Multiple	Unreliable/Imprecise fault
[Wei22]	SCA/ML	Multiple	Defeats 1th order masking
[FKK+22]	Rowhammer	Key Generation	KeyGen Failure Boosting
This work [HPP21]	Fault	Multiple	Unreliable fault

Allows for very unreliable fault, enlarges attack surface, only targets public data.

# Key Recovery from Decryption Failure Inequalities

Decryption failures leak information in form of inequalities:

- Exploited in wide variety of attacks including ours.
- Particularly hard to mitigate.
- Our attack strategy improved by [Del22], further enlarged attack surface.

Attacks require recovery method to solve for secret key:

- Partial information not considered; no security estimates.
- Error resistance increases #inequalities.
- General problem: combine belief propagation and algebraic methods?

Which techniques allow for key recovery from partially leaked decryption failure information?

# Key Recovery from Decryption Failure Inequalities

Decryption failures leak information in form of inequalities:

- Exploited in wide variety of attacks including ours.
- Particularly hard to mitigate.
- Our attack strategy improved by [Del22], further enlarged attack surface.

Attacks require recovery method to solve for secret key:

- Partial information not considered; no security estimates.
- Error resistance increases #inequalities.
- General problem: combine belief propagation and algebraic methods?

Which techniques allow for key recovery from partially leaked decryption failure information?



# Key Recovery from Decryption Failure Inequalities

Decryption failures leak information in form of inequalities:

- Exploited in wide variety of attacks including ours.
- Particularly hard to mitigate.
- Our attack strategy improved by [Del22], further enlarged attack surface.

Attacks require recovery method to solve for secret key:

- Partial information not considered; no security estimates.
- Error resistance increases #inequalities.
- General problem: combine belief propagation and algebraic methods?

**Which techniques allow for key recovery from partially leaked decryption failure information?**

# Security Estimates for Error-Tolerant Key Recovery

Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

Method	Inequalities	Error Resistant	Estimates
Pessl and Prokop [PP21]	8000	No	No
Hermelink et al. [HPP21]	5750	No	No
Delvaux [Del22]	9000	Yes	No
Dachman-Soled et al. [DDGR20, DGHK22]	$\geq 10000$	No	Yes

How can we combine the advantages of previous methods?

---

<sup>1</sup> Used for key recovery from such decryption failure information on widely available hardware in a concrete attack.

# Security Estimates for Error-Tolerant Key Recovery

Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

Method	Inequalities	Error Resistant	Estimates
Pessl and Prokop [PP21]	8000	No	No
Hermelink et al. [HPP21]	5750	No	No
Delvaux [Del22]	9000	Yes	No
Dachman-Soled et al. [DDGR20, DGHK22]	$\geq 10000$	No	Yes

How can we combine the advantages of previous methods?

---

<sup>1</sup> Used for key recovery from such decryption failure information on widely available hardware in a concrete attack.

# Security Estimates for Error-Tolerant Key Recovery

Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

Method	Inequalities	Error Resistant	Estimates
Pessl and Prokop [PP21]	8000	No	No
Hermelink et al. [HPP21]	5750	No	No
Delvaux [Del22]	9000	Yes	No
Dachman-Soled et al. [DDGR20, DGHK22]	$\geq 10000$	No	Yes

How can we combine the advantages of previous methods?

---

<sup>1</sup> Used for key recovery from such decryption failure information on widely available hardware in a concrete attack.

# Security Estimates for Error-Tolerant Key Recovery

Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

Method	Inequalities	Error Resistant	Estimates
Pessl and Prokop [PP21]	8000	No	No
Hermelink et al. [HPP21]	5750	No	No
Delvaux [Del22]	9000	Yes	No
Dachman-Soled et al. [DDGR20, DGHK22]	$\geq 10000$	No	Yes

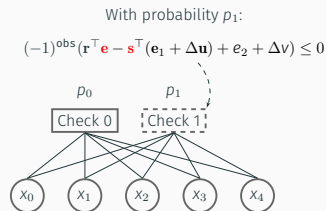
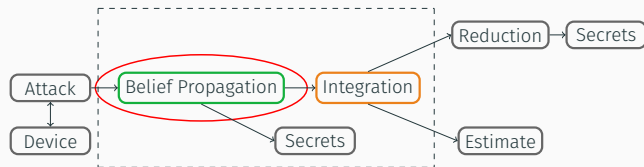
How can we combine the advantages of previous methods?

---

<sup>1</sup> Used for key recovery from such decryption failure information on widely available hardware in a concrete attack.

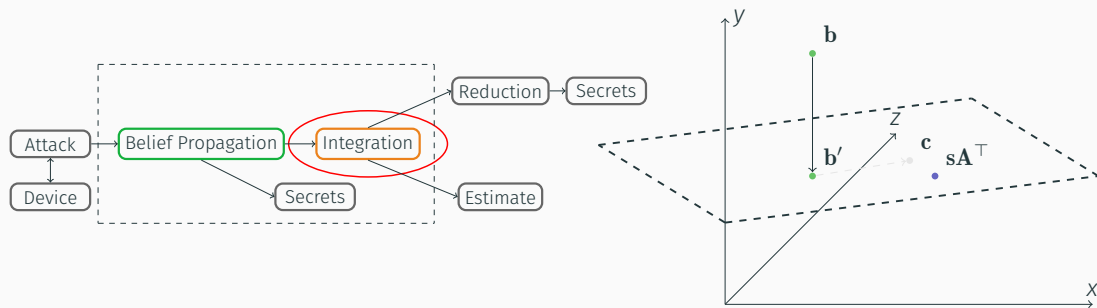
# Security Estimates for Error-Tolerant Key Recovery

New error-tolerant belief propagation with two-step lattice integration:



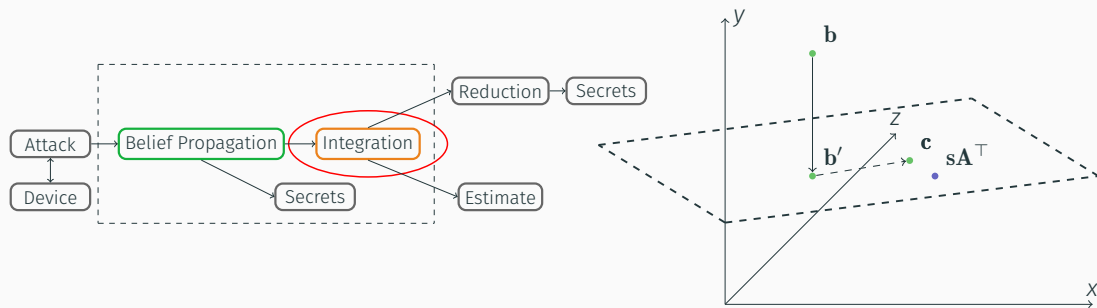
# Security Estimates for Error-Tolerant Key Recovery

New error-tolerant belief propagation with two-step lattice integration:



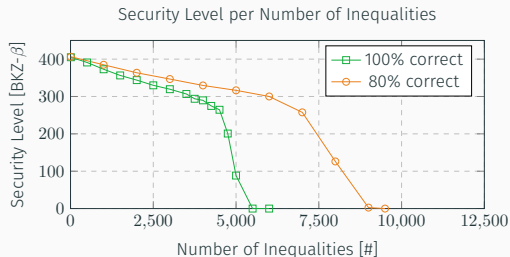
# Security Estimates for Error-Tolerant Key Recovery

New error-tolerant belief propagation with two-step lattice integration:





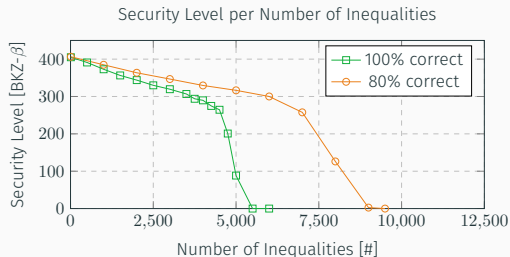
# Security Estimates for Error-Tolerant Key Recovery



Method	Inequalities	Error resistant	Security Estimates
[PP21]	8000	No	No
[HPP21]	5750	No	No
[Del22]	9000	Yes	No
[DDGR20, DGHK22]	$\geq 10000$	No	Yes
This work [HMS+23b]	5500	Yes	Yes

Hybrid approach: fewer inequalities, error-tolerant, and provides security estimates.

# Security Estimates for Error-Tolerant Key Recovery



Method	Inequalities	Error resistant	Security Estimates
[PP21]	8000	No	No
[HPP21]	5750	No	No
[Del22]	9000	Yes	No
[DDGR20, DGHK22]	$\geq 10000$	No	Yes
This work [HMS+23b]	5500	Yes	Yes

Hybrid approach: fewer inequalities, error-tolerant, and provides security estimates.

**NIST has standardized ML-KEM – a lattice-based scheme.**

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

NIST has standardized ML-KEM – a lattice-based scheme.

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

NIST has standardized ML-KEM – a lattice-based scheme.

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

NIST has standardized ML-KEM – a lattice-based scheme.

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

NIST has standardized ML-KEM – a lattice-based scheme.

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
- Statistical and algebraic tools relevant to a variety of attacks (used, e.g., in [DHP+22]).
- Extended assessment on vulnerabilities of major building blocks of modern lattice-based key encapsulation mechanisms.

# Publications

[HHP+21] Mike Hamburg, **Julius Hermelink**, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. “Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.4 (2021), pp. 88–113. URL: <https://doi.org/10.46586/tches.v2021.i4.88-113>

[HSST23] **Julius Hermelink**, Silvan Streit, Emanuele Strieder, and Katharina Thieme. “Adapting Belief Propagation to Counter Shuffling of NTTs”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.1 (2023), pp. 60–88. URL: <https://doi.org/10.46586/tches.v2023.i1.60-88>

[HMS+23b] **Julius Hermelink**, Erik Mårtensson, Simona Samardjiska, Peter Pessl, and Gabi Dreier Rodosek. “Belief Propagation Meets Lattice Reduction: Security Estimates for Error-Tolerant Key Recovery from Decryption Errors”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.4 (2023), pp. 287–317. URL: <https://doi.org/10.46586/tches.v2023.i4.287-317>

[HPS+20] **Julius Hermelink**, Thomas Pöppelmann, Marc Stöttinger, Yi Wang, and Yong Wan. “Quantum safe authenticated key exchange protocol for automotive application”. In: *18-th escar Europe : The World's Leading Automotive Cyber Security Conference (Konferenzveröffentlichung)*. 2020

[HPP21] **Julius Hermelink**, Peter Pessl, and Thomas Pöppelmann. “Fault-Enabled Chosen-Ciphertext Attacks on Kyber”. In: *Progress in Cryptology - INDOCRYPT 2021 - 22nd International Conference on Cryptology in India, Jaipur, India, December 12-15, 2021, Proceedings*. Ed. by Avishek Adhikari, Ralf Küsters, and Bart Preneel. Vol. 13143. Lecture Notes in Computer Science. Springer, 2021, pp. 311–334. URL: [https://doi.org/10.1007/978-3-030-92518-5\\_15](https://doi.org/10.1007/978-3-030-92518-5_15)



# References (1)

- [Aut23] Chromium Blog Authors. *Protecting Chrome Traffic with Hybrid Kyber KEM*. 2023. URL: <https://blog.chromium.org/2023/08/protecting-chrome-traffic-with-hybrid.html>.
- [BDH+21] Shivam Bhasin, Jan-Pieter D’Anvers, Daniel Heinz, Thomas Pöppelmann, and Michiel Van Beirendonck. “Attacking and Defending Masked Polynomial Comparison for Lattice-Based Cryptography”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.3 (2021), pp. 334–359. URL: <https://doi.org/10.46586/tches.v2021.i3.334-359>.
- [DDGR20] Dana Dachman-Soled, Léo Ducas, Huijing Gong, and Mélissa Rossi. “LWE with Side Information: Attacks and Concrete Security Estimation”. In: *Advances in Cryptology - CRYPTO 2020 - 40th Annual International Cryptology Conference, CRYPTO 2020, Santa Barbara, CA, USA, August 17-21, 2020, Proceedings, Part II*. Ed. by Daniele Micciancio and Thomas Ristenpart. Vol. 12171. Lecture Notes in Computer Science. Springer, 2020, pp. 329–358. URL: [https://doi.org/10.1007/978-3-030-56880-1\\_12](https://doi.org/10.1007/978-3-030-56880-1_12).
- [Del22] Jeroen Delvaux. “Roulette: A Diverse Family of Feasible Fault Attacks on Masked Kyber”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2022.4 (2022), pp. 637–660. URL: <https://doi.org/10.46586/tches.v2022.i4.637-660>.
- [DGHK22] Dana Dachman-Soled, Huijing Gong, Tom Hanson, and Hunter Kippen. “Revisiting Security Estimation for LWE with Hints from a Geometric Perspective”. In: *Advances in Cryptology - CRYPTO 2023 - 43rd Annual International Cryptology Conference, CRYPTO 2023, Santa Barbara, CA, USA, August 20-24, 2023, Proceedings, Part V*. Ed. by Helena Handschuh and Anna Lysyanskaya. Vol. 14085. Lecture Notes in Computer Science. Springer, 2022, pp. 748–781. URL: [https://doi.org/10.1007/978-3-031-38554-4\\_24](https://doi.org/10.1007/978-3-031-38554-4_24).

## References (2)

- [DHP+22] Jan-Pieter D’Anvers, Daniel Heinz, Peter Pessl, Michiel Van Beirendonck, and Ingrid Verbauwhede. “Higher-Order Masked Ciphertext Comparison for Lattice-Based Cryptography”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2022.2 (2022), pp. 115–139. URL: <https://doi.org/10.46586/tches.v2022.i2.115-139>.
- [EA24] Apple Security Engineering and Architecture. *iMessage with PQ3: The new state of the art in quantum-secure messaging at scale*. 2024. URL: <https://security.apple.com/blog/imessage-pq3/>.
- [FKK+22] Michael Fahr, Hunter Kippen, Andrew Kwong, Thinh Dang, Jacob Lichtinger, Dana Dachman-Soled, Daniel Genkin, Alexander Nelson, Ray A. Perlner, Arkady Yerukhimovich, and Daniel Apon. “When Frodo Flips: End-to-End Key Recovery on FrodoKEM via Rowhammer”. In: *Proceedings of the 2022 ACM SIGSAC Conference on Computer and Communications Security, CCS 2022, Los Angeles, CA, USA, November 7-11, 2022*. Ed. by Heng Yin, Angelos Stavrou, Cas Cremers, and Elaine Shi. ACM, 2022, pp. 979–993. URL: <https://doi.org/10.1145/3548606.3560673>.
- [GJN20] Qian Guo, Thomas Johansson, and Alexander Nilsson. “A Key-Recovery Timing Attack on Post-quantum Primitives Using the Fujisaki-Okamoto Transformation and Its Application on FrodoKEM”. In: *Advances in Cryptology - CRYPTO 2020 - 40th Annual International Cryptology Conference, CRYPTO 2020, Santa Barbara, CA, USA, August 17-21, 2020, Proceedings, Part II*. Ed. by Daniele Micciancio and Thomas Ristenpart. Vol. 12171. Lecture Notes in Computer Science. Springer, 2020, pp. 359–386. URL: [https://doi.org/10.1007/978-3-030-56880-1\\_13](https://doi.org/10.1007/978-3-030-56880-1_13).
- [HHP+21] Mike Hamburg, Julius Hermelink, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. “Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.4 (2021), pp. 88–113. URL: <https://doi.org/10.46586/tches.v2021.i4.88-113>.

## References (3)

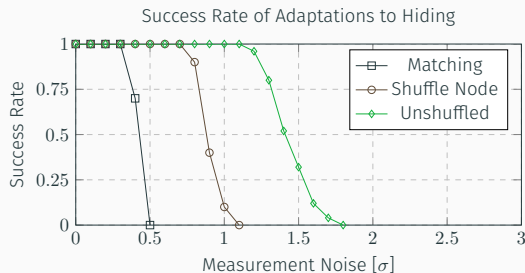
- [HMS+23b] Julius Hermelink, Erik Mårtensson, Simona Samardjiska, Peter Pessl, and Gabi Dreo Rodosek. “Belief Propagation Meets Lattice Reduction: Security Estimates for Error-Tolerant Key Recovery from Decryption Errors”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.4 (2023), pp. 287–317. URL: <https://doi.org/10.46586/tches.v2023.i4.287-317>.
- [HPP21] Julius Hermelink, Peter Pessl, and Thomas Pöppelmann. “Fault-Enabled Chosen-Ciphertext Attacks on Kyber”. In: *Progress in Cryptology - INDOCRYPT 2021 - 22nd International Conference on Cryptology in India, Jaipur, India, December 12-15, 2021, Proceedings*. Ed. by Avishek Adhikari, Ralf Küsters, and Bart Preneel. Vol. 13143. Lecture Notes in Computer Science. Springer, 2021, pp. 311–334. URL: [https://doi.org/10.1007/978-3-030-92518-5\\_15](https://doi.org/10.1007/978-3-030-92518-5_15).
- [HPS+20] Julius Hermelink, Thomas Pöppelmann, Marc Stöttinger, Yi Wang, and Yong Wan. “Quantum safe authenticated key exchange protocol for automotive application”. In: *18-th escar Europe : The World’s Leading Automotive Cyber Security Conference (Konferenzveröffentlichung)*. 2020.
- [HSST23] Julius Hermelink, Silvan Streit, Emanuele Strieder, and Katharina Thieme. “Adapting Belief Propagation to Counter Shuffling of NTTs”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.1 (2023), pp. 60–88. URL: <https://doi.org/10.46586/tches.v2023.i1.60-88>.
- [Kre23] Ehren Kret. *Quantum Resistance and the Signal Protocol*. 2023. URL: <https://signal.org/blog/pqxdh/>.

## References (4)

- [PP19] Peter Pessl and Robert Primas. “More Practical Single-Trace Attacks on the Number Theoretic Transform”. In: *Progress in Cryptology - LATINCRYPT 2019 - 6th International Conference on Cryptology and Information Security in Latin America, Santiago de Chile, Chile, October 2-4, 2019, Proceedings*. Ed. by Peter Schwabe and Nicolas Thériault. Vol. 11774. Lecture Notes in Computer Science. Springer, 2019, pp. 130–149. URL: [https://doi.org/10.1007/978-3-030-30530-7\\_7](https://doi.org/10.1007/978-3-030-30530-7_7).
- [PP21] Peter Pessl and Lukas Prokop. “Fault Attacks on CCA-secure Lattice KEMs”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.2 (2021), pp. 37–60. URL: <https://doi.org/10.46586/tches.v2021.i2.37-60>.
- [PPM17] Robert Primas, Peter Pessl, and Stefan Mangard. “Single-Trace Side-Channel Attacks on Masked Lattice-Based Encryption”. In: *Cryptographic Hardware and Embedded Systems - CHES 2017 - 19th International Conference, Taipei, Taiwan, September 25-28, 2017, Proceedings*. Ed. by Wieland Fischer and Naofumi Homma. Vol. 10529. Lecture Notes in Computer Science. Springer, 2017, pp. 513–533. URL: [https://doi.org/10.1007/978-3-319-66787-4\\_25](https://doi.org/10.1007/978-3-319-66787-4_25).
- [RPBC20] Prasanna Ravi, Romain Poussier, Shivam Bhasin, and Anupam Chattopadhyay. “On Configurable SCA Countermeasures Against Single Trace Attacks for the NTT - A Performance Evaluation Study over Kyber and Dilithium on the ARM Cortex-M4”. In: *Security, Privacy, and Applied Cryptography Engineering - 10th International Conference, SPACE 2020, Kolkata, India, December 17-21, 2020, Proceedings*. Ed. by Lejla Batina, Stjepan Picek, and Mainack Mondal. Vol. 12586. Lecture Notes in Computer Science. Springer, 2020, pp. 123–146. URL: [https://doi.org/10.1007/978-3-030-66626-2\\_7](https://doi.org/10.1007/978-3-030-66626-2_7).
- [Wei22] Andreas Weik. “Machine-Learning-based Side-Channel Attacks on Lattice-based Key Encapsulation Mechanisms”. Master’s Thesis at the Technical University of Munich 2022. Oct. 2022.

# Adapting Belief Propagation – Results

Real-world attacks have to take countermeasures into account.



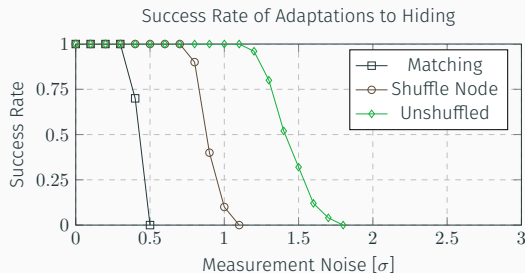
Our results show that

- Attacks not fully prevented by hiding countermeasures.
- However, noise tolerance reduced.
- Strongest form of shuffling protects requires vast computational resources.
- Large-scale adversaries might even circumvent coarse full shuffling.

Several hiding countermeasures can be circumvented; protected NTT still vulnerable.

# Adapting Belief Propagation – Results

Real-world attacks have to take countermeasures into account.



Our results show that

- Attacks not fully prevented by hiding countermeasures.
- However, noise tolerance reduced.
- Strongest form of shuffling protects requires vast computational resources.
- Large-scale adversaries might even circumvent coarse full shuffling.

Several hiding countermeasures can be circumvented; protected NTT still vulnerable.