

Side-Channel and Fault Attacks in Modern Lattice-Based Cryptography

Mündliche Prüfung – Julius Hermelink

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The Quantum Threat

Quantum computers threaten currently used asymmetric cryptography.



We have to assume that:

- Large-scale quantum computer break commonly used asymmetric schemes.
- Adversaries: harvest now, decrypt later.

Therefore, we need:

- Post-quantum asymmetric cryptography.
- Most pressing key exchanges.

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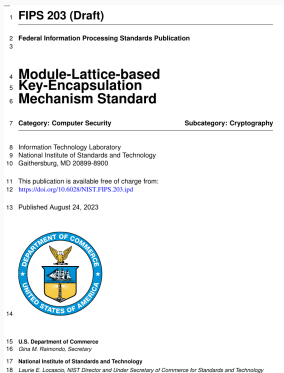
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The NIST Standardization Process

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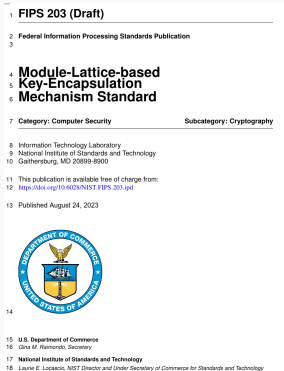


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- Four candidates already selected.
- Three are lattice-based.
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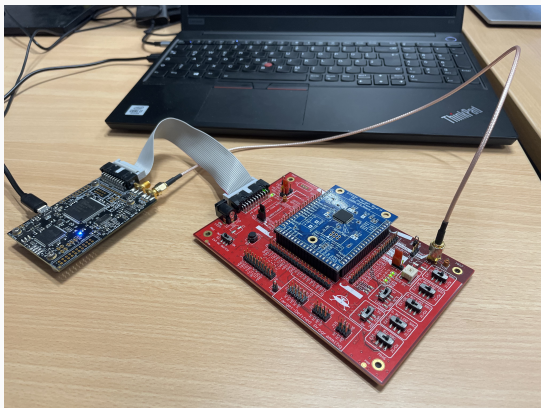
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ML-KEM used in Signal, Chrome, iMessage, ...



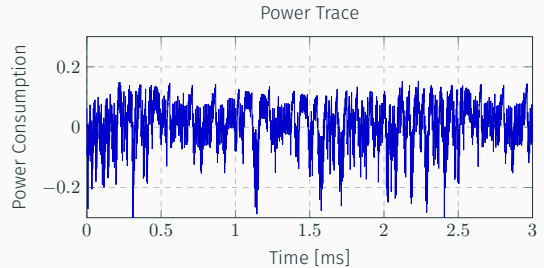
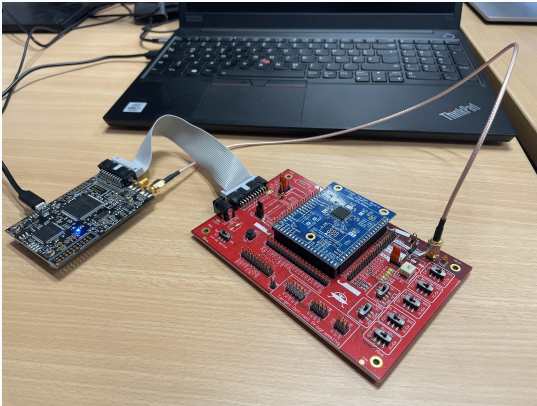
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Lattice-Based Cryptography uses different building blocks.

- Different underlying hard problems.
- Different multiplications (e.g., using number theoretic transforms).
- Error correction to recover message from noisy coefficients.
- Construction from PKE using FO-transforms to achieve IND-CCA security.
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Research Question 1

The number theoretic transform (NTT):

- Enables fast multiplication in several lattice-based schemes.
- Used at multiple points in all routines of ML-KEM.
- Inverse NTT processes data depending on the secret key during decryption.
- Previous work established (inverse) NTT as target for side-channel attacks.
- However, required noise levels limit attacks when targeting secret key.

RQ: To what extent is the number theoretic transform vulnerable to side-channel analysis.

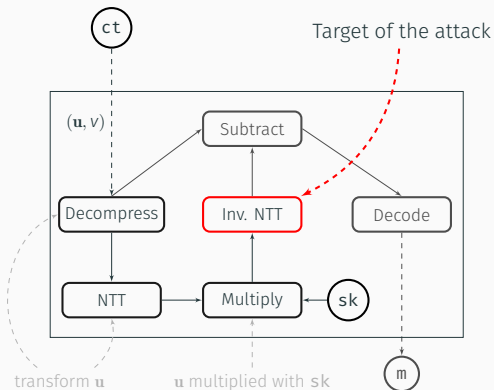
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Chosen-Ciphertext k-Trace Attacks – Idea



Previous work [PPM17, PP19]:

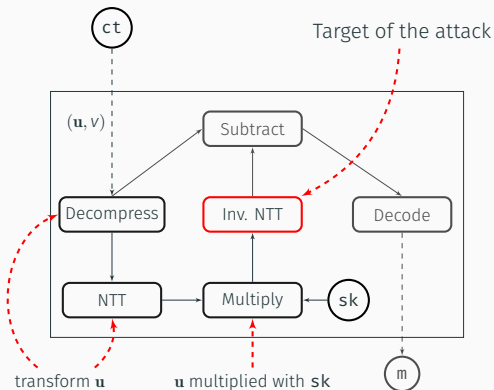
- Template attack on inv. NTT; then belief propagation.
- However, cannot target secret key with high noise tolerance.

Decryption as shown on the left:

- Ciphertext components are decompressed.
- Component is multiplied with secret.
- Results fed into the inv. number theoretic transform.

Attack strategy: Reduce entropy using compressible NTT-sparse chosen ciphertext.

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Our attack strategy for increased noise tolerance:

For targeted subkeys:

- Formulate as lattice problem.
- Run lattice reduction.
- Obtain compressible NTT-sparse $\mathbf{ct}$.

For each $\mathbf{ct}$:

- Record trace for $\mathbf{ct}$.
- Obtain distributions for intermediates.
- Run belief propagation; obtain subkeys.

Using the subkeys:

- Formulate key recovery using subkeys as lattice problem.
- Run lattice reduction.
- Obtain full key.

Lattice reduction is computationally expensive and slow but done offline.

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Adapting Belief Propagation to Counter Shuffling of NTTs

Real-world attacks have to take countermeasures into account.

Important classes of countermeasures:

- Masking circumvented.
- Hiding prevents these attacks.

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- Fine shuffling
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We propose two techniques against hiding:

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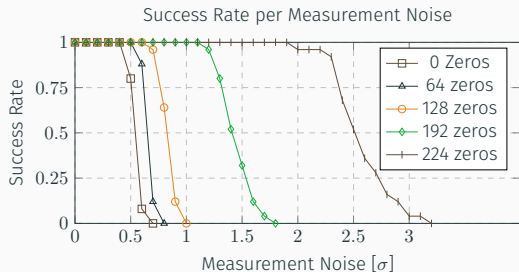
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Evaluation in the leakage models provided by previous work [PPM17, PP19].



Attack	Secret Key	Noise Tolerance	Traces	Hiding considered
[PPM17]	Yes	$\sigma \leq 0.6$	1	No
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This work	Yes	$\sigma \leq 1.7$ (3.1)	1-8	Yes

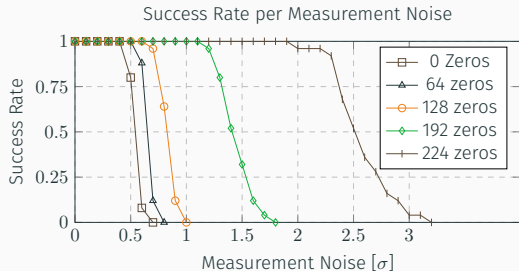
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Research Question 2

Learning with errors schemes recover the message from noisy coefficients:

- Observation whether introduced error causes failure: leaks information.
- Fujisaki-Okamoto (FO) transform achieves IND-CCA2 security.

Previous attacks:

- CCA to potentially cause failure; observe using SCA on comparison [GJN20, BDH⁺21].
- Or use fault against decoder to potentially cause failure and observe outcome [PP21].
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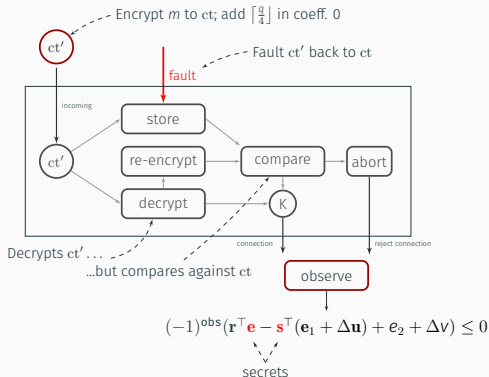
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Fault-Enabled Chosen-Ciphertext Attacks

Introduce error through chosen ciphertext, then correct with fault.

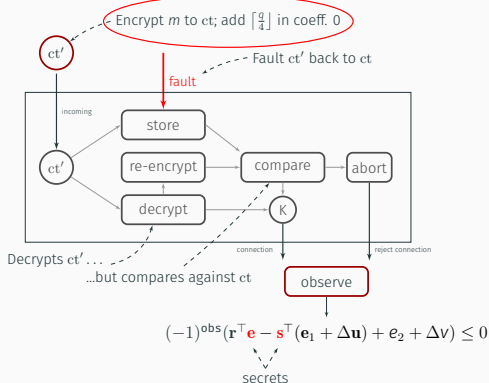
Our attack strategy:

- Ciphertext introduces error; fault corrects.
- Device decrypts ct' , but compares to ct .
- FO-comparison gives dec. failure oracle.
- Success can only occur if fault works.
- Allows for unreliable fault; attack surface over most of the execution time; may only target public data.



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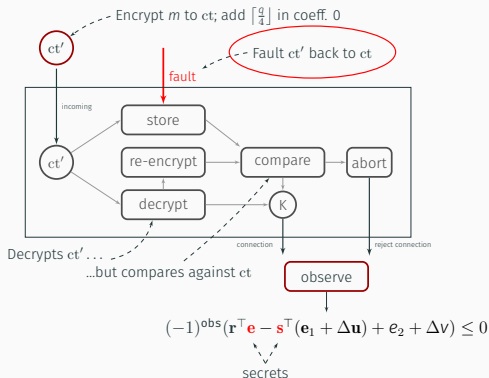


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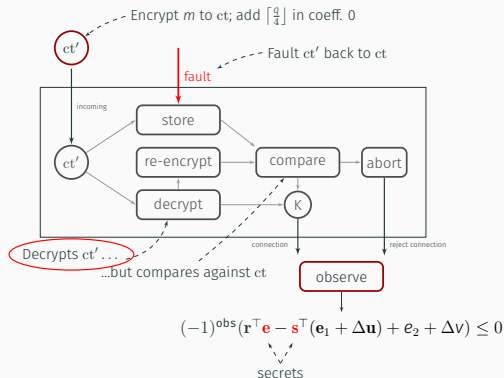


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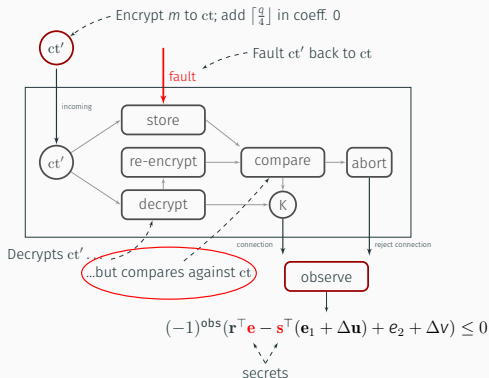
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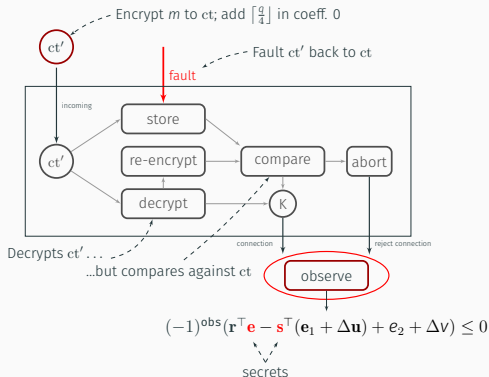


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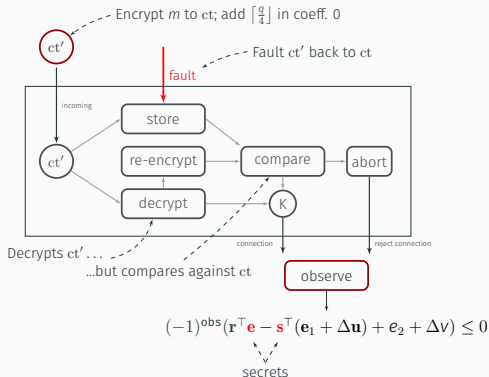


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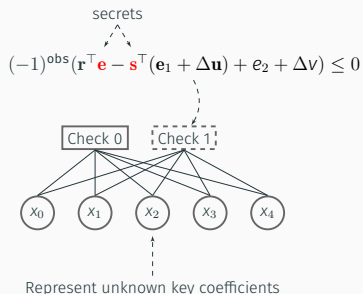
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Fault-Enabled Chosen-Ciphertext Attacks – BP

We propose solving decryption failure inequalities using belief propagation.



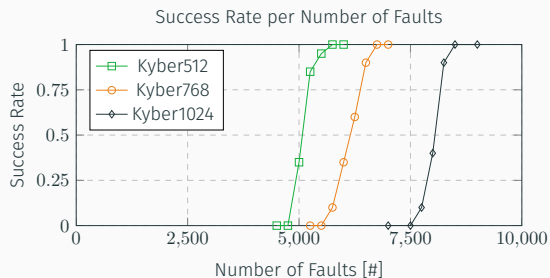
Belief propagation inspired by [PP21]:

- Check nodes represent inequalities.
- Variable nodes represent unknown coefficients.
- Priors are binomial distributions of secrets.

Requires fewer inequalities while being more computationally efficient.

Fault-Enabled Chosen-Ciphertext Attacks – Results

Gives a more general class of attacks resulting from our method.

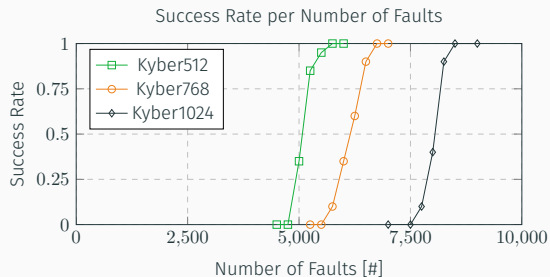


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- Particularly hard to mitigate.
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Attacks require recovery method to solve for secret key:

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Decryption failures in LWE leak information in form of inequalities.

Several methods to obtain secret from inequalities exist:

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Delvaux [Del22]	9000	Yes	No
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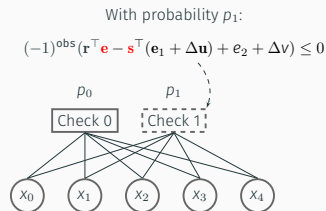
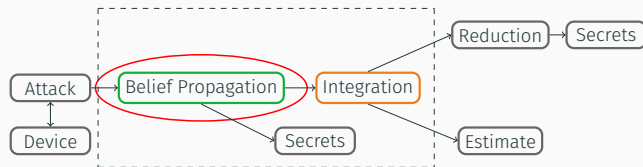
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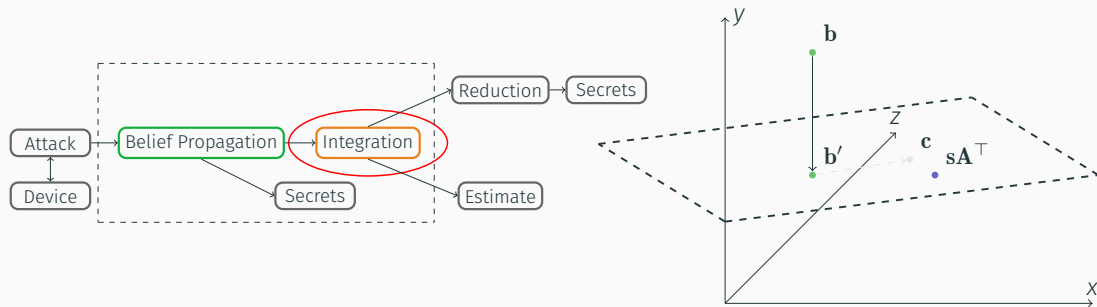
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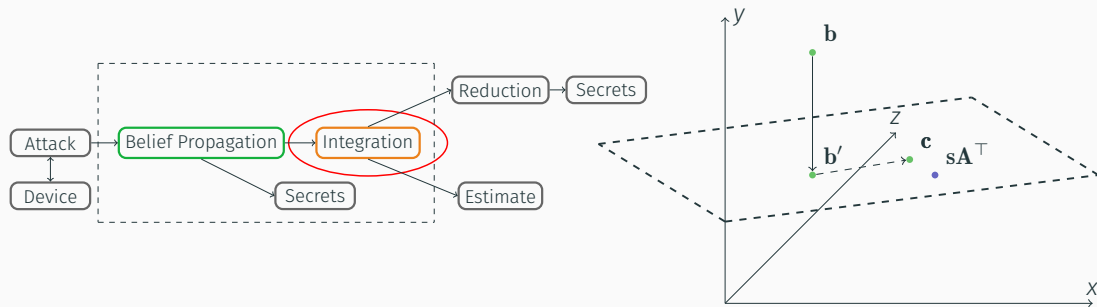
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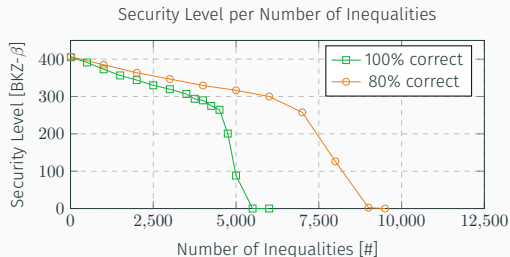


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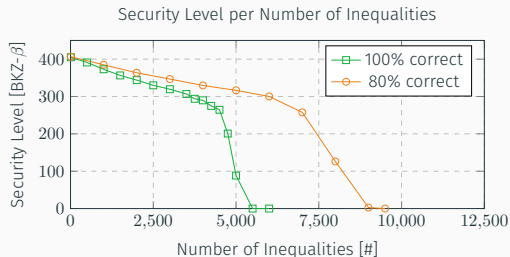
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- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
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NIST is standardizing ML-KEM – a lattice-based scheme.

Our work provides:

- Improvement on the state of the art in side-channel and fault attacks on lattice-based schemes.
- Attack strategies against lattice-based schemes enabling future attacks (e.g., [Del22]).
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Publications

[HHP+21] Mike Hamburg, **Julius Hermelink**, Robert Primas, Simona Samardjiska, Thomas Schamberger, Silvan Streit, Emanuele Strieder, and Christine van Vredendaal. “Chosen Ciphertext k-Trace Attacks on Masked CCA2 Secure Kyber”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2021.4 (2021), pp. 88–113. URL: <https://doi.org/10.46586/tches.v2021.i4.88-113>

[HSST23] **Julius Hermelink**, Silvan Streit, Emanuele Strieder, and Katharina Thieme. “Adapting Belief Propagation to Counter Shuffling of NTTs”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.1 (2023), pp. 60–88. URL: <https://doi.org/10.46586/tches.v2023.i1.60-88>

[HMS+23b] **Julius Hermelink**, Erik Mårtensson, Simona Samardjiska, Peter Pessl, and Gabi Dreier Rodosek. “Belief Propagation Meets Lattice Reduction: Security Estimates for Error-Tolerant Key Recovery from Decryption Errors”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.4 (2023), pp. 287–317. URL: <https://doi.org/10.46586/tches.v2023.i4.287-317>

[HPS+20] **Julius Hermelink**, Thomas Pöppelmann, Marc Stöttinger, Yi Wang, and Yong Wan. “Quantum safe authenticated key exchange protocol for automotive application”. In: *18-th escar Europe : The World's Leading Automotive Cyber Security Conference (Konferenzveröffentlichung)*. 2020

[HPP21] **Julius Hermelink**, Peter Pessl, and Thomas Pöppelmann. “Fault-Enabled Chosen-Ciphertext Attacks on Kyber”. In: *Progress in Cryptology - INDOCRYPT 2021 - 22nd International Conference on Cryptology in India, Jaipur, India, December 12-15, 2021, Proceedings*. Ed. by Avishek Adhikari, Ralf Küsters, and Bart Preneel. Vol. 13143. Lecture Notes in Computer Science. Springer, 2021, pp. 311–334. URL: https://doi.org/10.1007/978-3-030-92518-5_15

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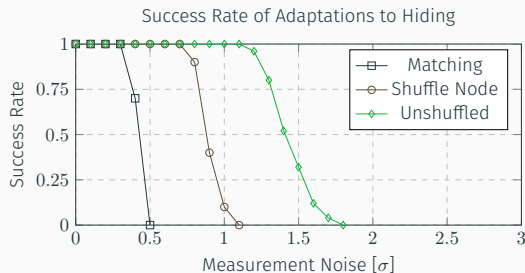
- [HMS+23b] Julius Hermelink, Erik Mårtensson, Simona Samardjiska, Peter Pessl, and Gabi Dreier Rodosek. “Belief Propagation Meets Lattice Reduction: Security Estimates for Error-Tolerant Key Recovery from Decryption Errors”. In: *IACR Trans. Cryptogr. Hardw. Embed. Syst.* 2023.4 (2023), pp. 287–317. URL: <https://doi.org/10.46586/tches.v2023.i4.287-317>.
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Adapting Belief Propagation – Results

Real-world attacks have to take countermeasures into account.



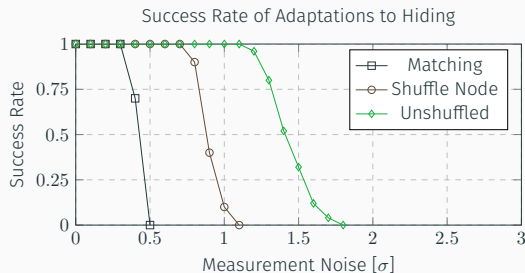
Our results show that

- Attacks not fully prevented by hiding countermeasures.
- However, noise tolerance reduced.
- Strongest form of shuffling protects requires vast computational resources.
- Large-scale adversaries might even circumvent coarse full shuffling.

Several hiding countermeasures can be circumvented; protected NTT still vulnerable.

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